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The Formal Inverse of the Period-Doubling Sequence

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Abstract

If p is a prime number, consider a p-automatic sequence $(u_n)_{n\geq 0}$, and let $U(X) = \sum_{n\geq 0} u_n X^n \in \mathbb{F}_p[[X]]$ be its generating function. Assume that there exists a formal power series $V(X) = \sum_{n\geq 0} v_n X^n \in \mathbb{F}_p[[X]]$ which is the compositional inverse of U, i.e., $U(V(X)) = X = V(U(X))$. The problem investigated in this paper is to study the properties of the sequence $(v_n)_{n\geq 0}$. The work was first initiated for the Thue-Morse sequence, and more recently the case of other sequences (variations of the Baum-Sweet sequence, variations of the Rudin-Shapiro sequence and generalized

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Thue-Morse sequences) has been treated. In this paper, we deal with the case of the period-doubling sequence. We first show that the sequence of indices at which the period-doubling sequence takes the value 0 (resp., 1) is not k-regular for any $k \geq 2$. Secondly, we give recurrence relations for its formal inverse, then we show that it is 2-automatic, and we also provide an automaton that generates it. Thirdly, we study the sequence of indices at which this formal inverse takes the value 1, and we show that it is not k-regular for any $k \geq 2$ by connecting it to the characteristic sequence of Fibonacci numbers. We leave as an open problem the case of the sequence of indices at which this formal inverse takes the value 0.

1 Introduction

Let us consider the following problem. Let p be a prime number. Let $\boldsymbol{u} = (u_n)_{n\geq 0}$ be a p-automatic sequence and let $U(X) = \sum_{n\geq 0} u_n X^n \in \mathbb{F}_p[[X]]$ be its generating function. Assume that there exists a formal power series $V(X) = \sum_{n\geq 0} v_n X^n \in \mathbb{F}_p[[X]]$ which is the compositional inverse of U, i.e., $U(V(X)) = X = V(U(X))$. What can be said about properties of the sequence $\boldsymbol{v} = (v_n)_{n\geq 0}$?

Gawron and Ulas [9] initiated work on this problem, and they consider the case where $\boldsymbol{u} = \boldsymbol{t}$ is the well-known Prouhet-Thue-Morse sequence. More precisely, they studied the sequence $\boldsymbol{c} = (c_n)_{n\geq 0}$, which is the sequence of coefficients of the compositional inverse of the generating function of the sequence $\boldsymbol{t}$. They called this sequence $\boldsymbol{c}$ the *inverse Prouhet-Thue-Morse sequence*. The 2-automaticity of $\boldsymbol{c}$ is easily deduced using Christol's theorem [5], but they also exhibited some recurrence relations satisfied by $\boldsymbol{c}$ and provided an automaton that generates $\boldsymbol{c}$. They studied two increasing sequences $\boldsymbol{a} = (a_n)_{n\geq 0}$ and $\boldsymbol{d} = (d_n)_{n\geq 0}$ respectively defined by

$$\{a_n \mid n \in \mathbb{N}\} = \{m \in \mathbb{N} \mid c_m = 1\},$$

and

$$\{d_n \mid n \in \mathbb{N}\} = \{m \in \mathbb{N} \mid c_m = 0\}.$$

In particular, they proved that $\boldsymbol{a}$ is 2-regular, but that $\boldsymbol{d}$ is not k-regular for any $k \geq 2$.

More recently, this work has been extended to other sequences [10, 11]. The author first obtains results similar to [9] for two variations of the Baum-Sweet sequence, and secondly for generalized Thue-Morse sequences (two specific cases) and two variations of the Rudin-Shapiro sequence.

In this paper, we consider the case where $\boldsymbol{u} = \boldsymbol{d}$ is the period-doubling sequence. This sequence is defined by $d_n := \nu_2(n+1) \bmod 2$, where the function ν_2 is the exponent of the highest power of 2 dividing its argument.

2 Background

In this section, we recall the necessary background for this paper; see, for instance, [4, 12, 13] for more details.

2.1 Combinatorics on words

Let A be a finite *alphabet*, i.e., a finite set consisting of *letters*. A *(finite) word* w over A is a finite sequence of letters belonging to A. If $w = w_n w_{n-1} \cdots w_0 \in A^*$ with $n \geq 0$ and $w_i \in A$ for all $i \in \{0, \ldots, n\}$, then the *length* $|w|$ of w is $n+1$, i.e., it is the number of letters that w contains. We let ε denote the empty word. This special word is the neutral element for concatenation of words, and its length is set to be 0. The set of all finite words over A is denoted by A^*, and we let $A^+ = A^* \setminus \{\varepsilon\}$ denote the set of non-empty finite words over A. For any $n \geq 0$, we let A^n denote the set of length-n words in A^*.

A finite word $w \in A^*$ is a *prefix* of another finite word $z \in A^*$ if there exists $u \in A^*$ such that $z = wu$. If A is ordered by $<$, the *lexicographic order* on A^*, which we denote by $<_{\text{lex}}$, is a total order on A^* induced by the order $<$ on the letters and defined as follows: $u <_{\text{lex}} v$ either if u is a strict prefix of v or if there exist $a, b \in A$ and $p \in A^*$ such that $a < b$, pa is a prefix of u and pb is a prefix of v.

If L is a subset of A^*, then L is called a *language* and its *complexity function* $\rho_L : \mathbb{N} \to \mathbb{N}$ is defined by $\rho_L(n) = L \cap A^n$.

An *infinite word* $\boldsymbol{w}$ over A is any infinite sequence over A. The set of all infinite words over A is denoted by A^ω. Note that in this paper infinite words are written in bold. To avoid any confusion, the infinite word $\boldsymbol{w} = w_0 w_1 w_2 \cdots$ will be written as $\boldsymbol{w} = w_0, w_1, w_2, \ldots$ if necessary.

If $\boldsymbol{w} \in A^\omega$, we define its *sequence of run lengths* to be an infinite sequence over $\mathbb{N} \cup \{\infty\}$ giving the number of adjacent identical letters. For example, the sequence of run lengths of $01^2 0^3 1^4 0^5 \cdots$ is $1, 2, 3, 4, 5, \ldots$.

A *morphism* on A is a map $\sigma : A^* \to A^*$ such that for all $u, v \in A^*$, we have $\sigma(uv) = \sigma(u)\sigma(v)$. In order to define a morphism, it suffices to provide the image of letters belonging to A. A morphism $\sigma : A^* \to A^*$ is *k-uniform* if $|\sigma(a)| = k$ for all $a \in A$. A 1-uniform morphism is called a *coding*. If there is a subalphabet $C \subset A$ such that $\sigma(C) \subset C^*$, then we call the restriction $\sigma_C := \sigma|_{C^*} : C^* \to C^*$ of σ to C a *submorphism* of σ.

A morphism $\sigma : A^* \to A^*$ is said to be *prolongable* on a letter $a \in A$ if $\sigma(a) = au$ with $u \in A^+$ and $\lim_{n \to +\infty} |\sigma^n(a)| = +\infty$. If σ is prolongable on a, then $\sigma^n(a)$ is a proper prefix of $\sigma^{n+1}(a)$ for all $n \geq 0$. Therefore, the sequence $(\sigma^n(a))_{n \geq 0}$ of finite words defines an infinite word $\boldsymbol{w}$ that is a fixed point of σ. In that case, the word $\boldsymbol{w}$ is called *pure morphic*. A *morphic* word is the morphic image of a pure morphic word.

Let M be a matrix with coefficients in $\mathbb{N}$. There exists permutation matrix P such that $P^{-1}MP$ is a upper block-triangular matrix with square blocks $M_1, \ldots, M_s$ on the main diagonal that are either irreducible matrices or zeroes. The *Perron-Frobenius* eigenvalue of M is $\max_{1 \leq i \leq s} \lambda_{M_i}$ where λ_{M_i} is the Perron-Frobenius eigenvalue of the matrix M_i.

Let $f : A^* \to A^*$ be a prolongable morphism having the infinite word $\boldsymbol{w}$ as a fixed point. Let α be the Perron-Frobenius eigenvalue of M_f. If all letters of A occur in $\boldsymbol{w}$, then $\boldsymbol{w}$ is said to be a *(pure) α-substitutive word*. If $g : A^* \to B^*$ is a coding, then $g(\boldsymbol{w})$ is said to be an *α-substitutive word*.

We say that two real numbers $\alpha, \beta > 1$ are *multiplicatively independent* if the only integers

k, ℓ such that $\alpha^k = \beta^\ell$ are $k = \ell = 0$. Otherwise, α and β are *multiplicatively dependent*. The following result can be found in [7].

Theorem 1 (Cobham-Durand)**.** *Let $\alpha, \beta > 1$ be two multiplicatively independent real numbers. Let $\boldsymbol{u}$ (resp., $\boldsymbol{v}$) be a pure α-substitutive (resp., pure β-substitutive) word. Let g and g' be two non-erasing morphisms. If $\boldsymbol{w} = g(\boldsymbol{u}) = g'(\boldsymbol{v})$, then $\boldsymbol{w}$ is ultimately periodic. In particular, if an infinite word is α-substitutive and β-substitutive, i.e., in the special case where g and g' are codings, then it is ultimately periodic.*

2.2 Abstract numeration systems, automatic sequences and regular sequences

An *abstract numeration system* (ANS) is a triple $S = (L, A, <)$ where L is an infinite regular language over a totally ordered alphabet $(A, <)$. The map $\mathrm{rep}_S : \mathbb{N} \to L$ is the one-to-one correspondence mapping $n \in \mathbb{N}$ onto the $(n+1)$st word in the genealogically ordered language L, which is called the *S-representation* of n. The S-representation of 0 is the first word in L. The inverse map is denoted by $\mathrm{val}_S : L \to \mathbb{N}$. If w is a word in L, $\mathrm{val}_S(w)$ is its *S-numerical value*. For instance, the base-k numeration system is an ANS; the Zeckendorff numeration system based on the Fibonacci numbers (with initial conditions 1 and 2) is also an ANS.

A *deterministic finite automaton with output* (DFAO) is a 6-tuple $\mathcal{A} = (Q, q_0, A, \delta, B, \mu)$, where Q is a finite set of *states*, $q_0 \in Q$ is the *initial state*, A is a finite *input alphabet*, $\delta : Q \times A \to Q$ is the *transition function*, B is a finite *output alphabet*, and $\mu : Q \to B$ is the *output function*. If $S = (L, A, <)$ is an ANS, we say that an infinite word $\boldsymbol{w} = w_0 w_1 w_2 \cdots \in B^{\mathbb{N}}$ is *S-automatic* if there exists a DFAO $\mathcal{A} = (Q, q_0, A, \delta, B, \mu)$ such that $x_n = \mu(\delta(q_0, \mathrm{rep}_S(n)))$ for all $n \geq 0$. The automaton $\mathcal{A}$ is called a *S-DFAO*.

When the *ANS* is the base-k numeration system with $k \geq 2$, we have the following theorem of Cobham [6].

Theorem 2 (Cobham's theorem on automatic sequences)**.** *An infinite word $\boldsymbol{w} \in B^{\mathbb{N}}$ is k-automatic if and only if there exist a k-uniform morphism $f : A^* \to A^*$ prolongable on a letter $a \in A$ and a coding $g : A^* \to B^*$ such that $\boldsymbol{w} = g(f^\omega(a))$.*

Let $\boldsymbol{u} = (u_n)_{n \geq 0}$ be an infinite sequence and let $k \geq 2$ be an integer. We define the *k-kernel* of $\boldsymbol{u}$ to be the set of subsequences

$$\mathcal{K}_k(\boldsymbol{u}) = \{(u_{k^i \cdot n + r})_{n \geq 0} \mid i \geq 0 \text{ and } 0 \leq r < k^i\}.$$

We say that a sequence $\boldsymbol{u}$ is *k-regular* if there exists a finite set S of sequences such that every sequence in $\mathcal{K}_k(\boldsymbol{u})$ is a $\mathbb{Z}$-linear combination of sequences of S. The following properties can be found in [4, 14].

Proposition 3. *Let $k \geq 2$ be an integer.*

(1) *If a sequence differs only in finitely many terms from a k-automatic sequence, then it is k-automatic.*

(2) *For all $m \geq 1$, a sequence is k-automatic if and only if it is k^m-automatic.*

(2) *If the integer sequence $(u_n)_{n\geq 0}$ is k-regular, then for all integers $m \geq 1$, the sequence $(u_n \bmod m)_{n\geq 0}$ is k-automatic.*

(3) *A sequence is k-regular and takes only finitely many values if and only if it is k-automatic.*

(4) *Let $(u_n)_{n\geq 0}$ be a k-regular sequence. Then for $a \geq 1$ and $b \geq 0$, the sequence $(u_{an+b})_{n\geq 0}$ is k-regular.*

(5) *Let $\boldsymbol{u} = (u_n)_{n\geq 0}$ be a sequence, and let $\boldsymbol{v} = (u_{n+1} - u_n)_{n\geq 0}$ be the first difference of $\boldsymbol{u}$. Then $\boldsymbol{u}$ is k-regular if and only if $\boldsymbol{v}$ is k-regular.*

2.3 Formal power series

Let $k \geq 2$. The *ring $\mathbb{F}_k[[X]]$ of formal power series with coefficients in the field $\mathbb{F}_k = \{0, 1, \ldots, k-1\}$* is defined by

$$\mathbb{F}_k[[X]] = \left\{ \sum_{n\geq 0} a_n X^n \mid a_n \in \mathbb{F}_k \right\}.$$

We let $\mathbb{F}_k(X)$ denote the *the field of rational functions.* We say that a formal series $A(X) = \sum_{n\geq 0} a_n X^n$ is *algebraic (over $\mathbb{F}_k(X)$)* if there exist an integer $d \geq 1$ and polynomials $P_0(X)$, $P_1(X)$, ..., $P_d(X)$, with coefficients in $\mathbb{F}_k$ and not all zero, such that

$$P_0 + P_1 A + P_2 A^2 + \cdots + P_d A^d = 0.$$

With an infinite sequence $\boldsymbol{w} = (w_n)_{n\in\mathbb{N}}$ over $\{0, 1, \ldots, k-1\}$, we can associate a formal series $W(X) = \sum_{n\geq 0} w_n X^n$ over $\mathbb{F}_k[[X]]$, which is called the *generating function* of $\boldsymbol{w}$. In the case where $k = p$ is a prime number, and if $w_0 = 0$ and w_1 is invertible in $\mathbb{F}_p$, then the series $W(X)$ is *invertible* in $\mathbb{F}_p[[X]]$, i.e., there exists a series $U(X) \in \mathbb{F}_p[[X]]$ such that $W(U(X)) = X = U(W(X))$. The formal series $U(X)$ is called the *(formal) inverse* of $W(X)$.

3 The period-doubling sequence

The following definition can be found in [4].

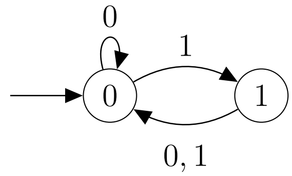

Figure 1: The 2-DFAO generating the period-doubling sequence $\boldsymbol{d}$.

Definition 4. Consider the period-doubling sequence (indexed by `A096268` in [15])

$$\boldsymbol{d} = (d_n)_{n\geq 0} = 0100010101000100010001000\cdots.$$

This sequence is defined by $d_n := \nu_2(n+1) \bmod 2$, where the function ν_2 is the exponent of the highest power of 2 dividing its argument. Alternatively, we have $\boldsymbol{d} = h^\omega(0)$, where $h(0) = 01$ and $h(1) = 00$. Since h is a 2-uniform morphism, then the period doubling sequence $\boldsymbol{d}$ is 2-automatic. The 2-DFAO drawn in Figure 1 generates the period-doubling sequence $\boldsymbol{d}$. Note that this automaton reads its input from most significant digit to least significant digit.

Let us define two increasing sequences $\boldsymbol{o} = (o_n)_{n\geq 0}$ and $\boldsymbol{z} = (z_n)_{n\geq 0}$ respectively satisfying $\{o_n \mid n \in N\} = \{m \in N \mid d_m = 1\}$ and $\{z_n \mid n \in N\} = \{m \in N \mid d_m = 0\}$. We have

$$\begin{aligned}\boldsymbol{o} &= 1, 5, 7, 9, 13, 17, 21, 23, 25, 29, 31, 33, 37, 39, 41, 45, 49, 53, 55, 57, 61, 65, 69, 71, 73, 77, \ldots,\\ \boldsymbol{z} &= 0, 2, 3, 4, 6, 8, 10, 11, 12, 14, 15, 16, 18, 19, 20, 22, 24, 26, 27, 28, 30, 32, 34, 35, 36, 38, 40, \ldots.\end{aligned}$$

Those two sequences are indexed by `A079523` and `A121539` in [15]. Observe that the binary expansions of the terms of $\boldsymbol{o}$ (resp., $\boldsymbol{z}$) end with an odd (resp., even) number of 1's. This can be seen if one considers the language accepted by the 2-DFAO in Figure 1 where the final state is the one outputting 1 (resp., 0). In the following, we study the regularity of the sequences $\boldsymbol{o}$ and $\boldsymbol{z}$.

Proposition 5. *The sequence $\boldsymbol{z} = (z_n)_{n\geq 0}$ is not k-regular for any $k \in \mathbb{N}_{\geq 2}$.*

Proof. Let $\bar{\boldsymbol{d}}$ be the image of $\boldsymbol{d}$ under the exchange morphism $E : \{0,1\}^* \to \{0,1\}^* : 0 \mapsto 1, 1 \mapsto 0$. In particular, $\bar{\boldsymbol{d}}$ is the fixed point of the morphism $h'(0) = 11$ and $h'(1) = 10$ starting with 1. We also have

$$\boldsymbol{z} = \{m \in \mathbb{N} \mid d_m = 0\} = \{m \in \mathbb{N} \mid \bar{d}_m = 1\}.$$

The sequence $\bar{\boldsymbol{d}}$ is related to the Thue-Morse sequence in the following way. Let $\boldsymbol{t} = (t_n)_{n\geq 0}$ be the Thue-Morse sequence, i.e., the fixed point of the morphism $\tau : \{0,1\}^* \to \{0,1\}^* : 0 \mapsto 01, 1 \mapsto 10$ which starts with 0. In fact, the sequence $\bar{\boldsymbol{d}}$ is the first difference modulo 2 of the Thue-Morse sequence $\boldsymbol{t}$ [3], i.e., $\bar{\boldsymbol{d}} = (t_{n+1} - t_n \bmod 2)_{n\geq 0}$.

In other words, the sequence $\boldsymbol{z}$ of positions of 1's in $\bar{\boldsymbol{d}}$ is exactly the sequence of positions in the Thue-Morse sequence $\boldsymbol{t}$ where the letters 0 and 1 alternate. Consequently, the first difference of $\boldsymbol{z}$, which is the first difference between the positions of 1's in $\bar{\boldsymbol{d}}$, gives the length of the blocks of consecutive identical letters in $\boldsymbol{t}$, i.e., it is the sequence of run lengths of $\boldsymbol{t}$.

However, the sequence of run lengths of $\boldsymbol{t}$ is the sequence $\boldsymbol{p} = (p_n)_{n\geq 0}$ which is the fixed point of the morphism $f : \{1,2\}^* \to \{1,2\}^* : 1 \mapsto 121, 2 \mapsto 12221$ which starts with 1 [2]. This sequence $\boldsymbol{p}$ is not 2-automatic [1], and by Proposition 3, $\boldsymbol{p}$ is not 2^m-automatic for any $m \geq 1$. Let us show that $\boldsymbol{p}$ is not k-automatic for any integer $k \geq 2$. Suppose that $\boldsymbol{p}$ is k-automatic for some integer $k \geq 2$ which is not a power of 2. Then, by Theorem 2, $\boldsymbol{p}$ is the image under a coding of the fixed point of a k-uniform morphism whose Perron-Frobenius eigenvalue is k. Since the Perron-Frobenius eigenvalue of f is 2, then by Theorem 1, $\boldsymbol{p}$ is ultimately periodic, which is impossible.

Now since $\boldsymbol{p}$ takes only two different values, $\boldsymbol{p}$ is not k-regular for any $k \geq 2$ by Proposition 3. Since $\boldsymbol{p}$ is the first difference of $\boldsymbol{z}$, then $\boldsymbol{z}$ is not k-regular for any $k \geq 2$ again by Proposition 3. $\square$

The next lemma gives two other morphisms that generate the period-doubling sequence $\boldsymbol{d}$. Those morphisms are helpful to locate the positions of 1's in $\boldsymbol{d}$.

Lemma 6. *Let $f : \{2,4\}^* \to \{2,4\}^* : 2 \mapsto 242, 4 \mapsto 24442$ and $g : \{2,4\}^* \to \{0,1\}^* : 2 \mapsto 01, 4 \mapsto 0001$. For all $n \geq 1$, we have $h^{2n+1}(0) = g(f^n(2))$ and $h^{2n+1}(10) = g(f^n(4))$. In particular, $\boldsymbol{d} = h^\omega(0) = g(f^\omega(2))$.*

Proof. We proceed by induction on $n \geq 1$. The case $n = 1$ can easily be checked by hand. Now assume that $n \geq 1$ and suppose that the result holds for all $m \leq n$. We have

$$h^{2(n+1)+1}(0) = h^{2n+1}(0100) = h^{2n+1}(0)h^{2n+1}(10)h^{2n+1}(0).$$

Now, by induction hypothesis, we find

$$h^{2(n+1)+1}(0) = g(f^n(2))g(f^n(4))g(f^n(2)) = g(f^n(242)) = g(f^{n+1}(2)),$$

as expected. Similarly, we have

$$h^{2(n+1)+1}(10) = h^{2n+1}(01010100) = h^{2n+1}(0)h^{2n+1}(10)h^{2n+1}(10)h^{2n+1}(10)h^{2n+1}(0),$$

and by induction hypothesis, we get

$$h^{2(n+1)+1}(0) = g(f^n(2))g(f^n(4))g(f^n(4))g(f^n(4))g(f^n(2)) = g(f^n(24442)) = g(f^{n+1}(4)).$$

The particular case can be deduced from the first equality of the statement. $\square$

Proposition 7. *The sequence $\boldsymbol{o} = (o_n)_{n\geq 0}$ is not k-regular for any $k \in \mathbb{N}_{\geq 2}$.*

Proof. By Lemma 6, we know that $\boldsymbol{d} = g(f^\omega(2))$ with $f : \{2,4\}^* \to \{2,4\}^* : 2 \mapsto 242, 4 \mapsto 24442$ and $g : \{2,4\}^* \to \{0,1\}^* : 2 \mapsto 01, 4 \mapsto 0001$. Observe that $|g(2)| = 2$ and $|g(4)| = 4$, and the letter 1 occurs only once at the end of $g(2)$ (resp., $g(4)$). Consequently, the first difference of the positions of 1's in $\boldsymbol{d}$ — which is the first difference of $\boldsymbol{o}$ — is given by the shift of the sequence $f^\omega(2)$, i.e., we drop the first term. By the proof of Proposition 5, we know that $f^\omega(2)$ is not k-regular for any $k \geq 2$. By Proposition 3, $\boldsymbol{o}$ is not k-regular for any $k \geq 2$. □

Remark 8. Using an argument similar to the one of the proof of Proposition 7, one can also get another way of proving Proposition 5.

4 The formal inverse of the period-doubling word

Let $D(X) = \sum_{n\geq 0} d_n X^n$ be the generating function of the period-doubling sequence $\boldsymbol{d}$. Since $d_0 = 0$ and $d_1 = 1$ is invertible in $\mathbb{F}_2$, then the series $D(X)$ is invertible in $\mathbb{F}_2[[X]]$, i.e., there exists a series

$$U(X) = \sum_{n\geq 0} u_n X^n \in \mathbb{F}_2[[X]]$$

such that $D(U(X)) = X = U(D(X))$. We want to describe the sequence $\boldsymbol{u} = (u_n)_{n\geq 0}$. Mimicking [9], the first step is to get recurrence relations for the coefficients $(u_n)_{n\geq 0}$ of the series $U(X)$. To that aim, recall the following result; see [5, p. 412].

Lemma 9. *The generating function $D(X) = \sum_{n\geq 0} d_n X^n$ of the period-doubling sequence $\boldsymbol{d}$ satisfies*

$$X(1+X^2)D(X)^2 + (1+X^2)D(X) + X = 0$$

over $\mathbb{F}_2[[X]]$.

Proof. Observe that, since $\boldsymbol{d} = h^\omega(0)$, we have $d_{2n} = 0$ and $d_{2n+1} = 1 - d_n$ for all $n \geq 0$. Thus we have

$$D(X) = \sum_{n\geq 0} d_n X^n = \sum_{n\geq 0} d_{2n} X^{2n} + \sum_{n\geq 0} d_{2n+1} X^{2n+1} = X \sum_{n\geq 0} X^{2n} - X \sum_{n\geq 0} d_n X^{2n}.$$

Now recall that, for any prime p and for any series $F(X)$ in $\mathbb{F}_p[[X]]$, we have $1/(1-X) = \sum_{n\geq 0} X^n$. Consequently,

$$D(X) = \frac{X}{1-X^2} - XD(X^2).$$

Now working over $\mathbb{F}_2[[X]]$, we have

$$X(1+X^2)D(X^2) + (1+X^2)D(X) + X = 0,$$

and since for any prime p and for any series $F(X)$ in $\mathbb{F}_p[[X]]$, we have $F(X)^p = F(X^p)$, we find

$$X(1+X^2)D(X)^2 + (1+X^2)D(X) + X = 0,$$

as desired. □

To prove the next result, we follow the method from [9].

Proposition 10. *The series* $U(X) = \sum_{n\geq 0} u_n X^n$ *satisfies each of the following polynomial equations*

$$\begin{aligned} X^2U(X)^3 + XU(X)^2 + (X^2+1)U(X) + X = 0,\\ X^3U(X)^4 + X^3U(X)^2 + U(X) + X = 0 \end{aligned}$$

over $\mathbb{F}_2[[X]]$*. In particular, the sequence* $\boldsymbol{u} = (u_n)_{n\geq 0}$ *verifies* $u_0 = 0$*,* $u_1 = 1$*, and over* $\mathbb{F}_2$

$$\begin{cases} u_{2n} = 0 \quad \forall\, n \geq 0, \\ u_{4n+1} = u_{2n-1} \quad \forall\, n \geq 1, \\ u_{4n+3} = u_n \quad \forall\, n \geq 0. \end{cases}$$

Proof. First, let us rewrite the equation from Lemma 9 in terms of X. We get

$$D(X)^2X^3 + D(X)X^2 + (D(X)^2+1)X + D(X) = 0.$$

In this new equation, replace X by $U(X)$ to obtain

$$D(U(X))^2U(X)^3 + D(U(X))U(X)^2 + (D(U(X))^2+1)U(X) + D(U(X)) = 0.$$

Since $U(X)$ is the formal inverse of $D(X)$, we actually have

$$X^2U(X)^3 + XU(X)^2 + (X^2+1)U(X) + X = 0, \tag{1}$$

which is the first equation of the statement. This in turn implies that, over $\mathbb{F}_2[[X]]$,

$$U(X)^3 = \frac{XU(X)^2 + (X^2+1)U(X) + X}{X^2}. \tag{2}$$

Now multiply (1) by $U(X)$ and replace $U(X)^3$ by its value (2). We obtain first

$$X^2U(X)^4 + XU(X)^3 + (X^2+1)U(X)^2 + XU(X) = 0,$$

and so

$$\begin{aligned} & X^2U(X)^4 + X\left(\frac{XU(X)^2 + (X^2+1)U(X) + X}{X^2}\right) + (X^2+1)U(X)^2 + XU(X) = 0 \\ & \Rightarrow X^3U(X)^4 + XU(X)^2 + (X^2+1)U(X) + X + (X^3+X)U(X)^2 + X^2U(X) = 0 \\ & \Rightarrow X^3U(X)^4 + (X^3+2X)U(X)^2 + (2X^2+1)U(X) + X = 0. \end{aligned}$$

Working over $\mathbb{F}_2[[X]]$, this equality becomes

$$X^3U(X)^4 + X^3U(X)^2 + U(X) + X = 0 \Leftrightarrow X^3U(X^4) + X^3U(X^2) + U(X) + X = 0,$$

which is the second equation of the statement.

Let us now prove that the recurrence relations for the sequence $\boldsymbol{u}$ hold. Writing $U(X) = \sum_{n\geq 0} u_n X^n$ in the second equation proven above, we find

$$\begin{aligned} & X^3 \sum_{n\geq 0} u_n X^{4n} + X^3 \sum_{n\geq 0} u_n X^{2n} + \sum_{n\geq 0} u_n X^n + X = 0 \\ \Leftrightarrow & \sum_{n\geq 0} u_n X^{4n+3} + \sum_{n\geq 0} u_n X^{2n+3} + \sum_{n\geq 0} u_n X^n + X = 0. \end{aligned}$$

Let us inspect the coefficients in the last equality. We immediately have $u_0 = 0$ and $u_1 = 1$ over $\mathbb{F}_2$. Since the exponents $4n+3$ and $2n+3$ are odd for all $n \geq 0$, we also get that, over $\mathbb{F}_2$,

$$u_{2n} = 0 \quad \forall\, n \geq 0.$$

Looking at the coefficient of X^{4n+3}, we obtain

$$u_n + u_{2n} + u_{4n+3} = 0 \quad \forall\, n \geq 0,$$

which implies that $u_{4n+3} = u_n$ over $\mathbb{F}_2$ for all $n \geq 0$. Let us now find the coefficient of X^{4n+1} for $n \geq 1$. We have

$$u_{2n-1} + u_{4n+1} = 0 \quad \forall\, n \geq 1,$$

giving $u_{4n+1} = u_{2n-1}$ over $\mathbb{F}_2$ for all $n \geq 1$. As a consequence, the sequence $\boldsymbol{u} = (u_n)_{n\geq 0}$ verifies $u_0 = 0$, $u_1 = 1$, and satisfies the following recurrence relations over $\mathbb{F}_2$

$$\begin{cases} u_{2n} = 0 \quad \forall\, n \geq 0, \\ u_{4n+1} = u_{2n-1} \quad \forall\, n \geq 1, \\ u_{4n+3} = u_n \quad \forall\, n \geq 0. \end{cases}$$

□

From now on, the sequence $\boldsymbol{u} = (u_n)_{n\geq 0}$ will be referred to as the *inverse period-doubling sequence*, iPD sequence for short (sequence `A317542` in [15]). We have

$$\boldsymbol{u} = (u_n)_{n\geq 0} = 0100010100000100010000010000010100000 1000 \cdots .$$

Remark 11. We have $d_n = u_n$ for all $n \leq 8$, but observe that

$$1 = d_{4\cdot 2+1} = d_9 \neq u_9 = u_{4\cdot 2+1} u = u_{2\cdot 2-1} = u_3 = 0.$$

In the following, we show that $\boldsymbol{u}$ is 2-automatic, and we also provide an automaton that generates $\boldsymbol{u}$, which is deduced from the recurrence relations in Lemma 13.

Corollary 12. *The sequence $\boldsymbol{u} = (u_n)_{n\geq 0}$ is 2-automatic.*

Proof. From Proposition 10, it follows that the formal power series $U(X)$ is algebraic over $\mathbb{F}_2(X)$. By Christol's theorem, the sequence $\boldsymbol{u}$ is thus 2-automatic. □

Lemma 13. *For all* $n \geq 0$, $r_1 \in \{0, 2\}$, $r_2 \in \{0, 2, 4, 6\}$ *and* $r_3 \in \{0, 2, 4, 6, 8, 10, 12, 14\}$, *we have*

$$u_n = u_{4n+3} = u_{16n+15}, \tag{3}$$
$$u_{2n} = u_{4n+r_1} = u_{8n+r_2} = u_{8n+3} = u_{16n+r_3} = u_{16n+3} = u_{16n+9} = u_{16n+11} = 0, \tag{4}$$
$$u_{2n+1} = u_{8n+7}, \tag{5}$$
$$u_{4n+1} = u_{8n+5} = u_{16n+1} = u_{16n+7} = u_{16n+13}, \tag{6}$$
$$u_{8n+1} = u_{16n+5}. \tag{7}$$

Proof. We make an extensive use of the recurrence relations from Proposition 10. We show that the 2-kernel $\mathcal{K}_2(\boldsymbol{u})$ is finitely generated by the sequences $(u_n)_{n\geq0}$, $(u_{2n})_{n\geq0}$, $(u_{2n+1})_{n\geq0}$, $(u_{4n+1})_{n\geq0}$ and $(u_{8n+1})_{n\geq0}$.

The first equality in (3) is directly given by Proposition 10. For all $n \geq 0$, we have

$$u_{16n+15} = u_{4(4n+3)+3} = u_{4n+3} = u_n$$

using Proposition 10 twice since $n, 4n + 3 \geq 0$.

Let us show (4). From Proposition 10, it is clear that for all $n \geq 0$,

$$u_{2n} = 0 = u_{4n+r_1} = u_{8n+r_2} = u_{16n+r_3}.$$

Now for all $n \geq 0$, we have

$$u_{8n+3} = u_{4(2n)+3} = u_{2n} = 0,$$
$$u_{16n+3} = u_{4(4n)+3} = u_{4n} = u_{2n} = 0,$$

and

$$u_{16n+11} = u_{4(4n+2)+3} = u_{4n+2} = u_{2n} = 0,$$

using Proposition 10 since $2n, 4n, 4n + 2 \geq 0$. Similarly, for all $n \geq 0$, we have $4n + 2 \geq 1$, thus Proposition 10 gives

$$u_{16n+9} = u_{4(4n+2)+1} = u_{2(4n+2)-1} = u_{8n+3} = u_{2n} = 0,$$

where the next-to-last equality comes from (4) above.

Let us prove (5). For all $n \geq 0$, we have

$$u_{8n+7} = u_{4(2n+1)+3} = u_{2n+1},$$

using Proposition 10 since $2n + 1 \geq 0$.

Let us show that (6) holds. For all $n \geq 0$, we have

$$u_{8n+5} = u_{4(2n+1)+1} = u_{2(2n+1)-1} = u_{4n+1},$$

$$u_{16n+7} = u_{4(4n+1)+3} = u_{4n+1},$$

and

$$u_{16n+13} = u_{4(4n+3)+1} = u_{2(4n+3)-1} = u_{8n+5} = u_{4n+1},$$

using Proposition 10 since $2n+1, 4n+3 \geq 1$ and $4n+1 \geq 0$. Now we prove that $u_{16n+1} = u_{4n+1}$ for all $n \geq 0$. The result is trivial when $n = 0$ for we have $u_{16n+1} = u_1 = u_{4n+1}$. Now suppose that $n \geq 1$. We first obtain from Proposition 10 that

$$u_{16n+1} = u_{4(4n)n+1} = u_{2(4n)-1} = u_{8n-1}.$$

Writing $n = m + 1$ with $m \geq 0$, we then get

$$u_{16n+1} = u_{8n-1} = u_{8m+7} = u_{2m+1}$$

where the last equality comes from (5) since $m \geq 0$. Consequently,

$$u_{16n+1} = u_{2m+1} = u_{2(m+1)-1} = u_{2n-1} = u_{4n+1}$$

using Proposition 10 for the last equality since $n \geq 1$. This gives the expected recurrence relation.

Finally, for all $n \geq 0$, we have $4n + 1 \geq 0$, so Proposition 10 implies that

$$u_{16n+5} = u_{4(4n+1)+1} = u_{2(4n+1)-1} = u_{8n+1},$$

which proves (7). □

Using the recurrence relations from Lemma 13, we build the 2-DFAO in Figure 2, which generates the iPD sequence $\boldsymbol{u}$. Note that this automaton reads its input from least significant digit to most significant digit.

Since the iPD sequence $\boldsymbol{u}$ takes the values 0 and 1, it can also be considered as a sequence of complex numbers. We now obtain the transcendence of its generating function.

Proposition 14. *The formal power series $U(X) = \sum_{n\geq 0} u_n X^n \in \mathbb{C}[[X]]$ is transcendental over $\mathbb{C}(X)$.*

Proof. A classical result of Fatou states that a power series whose coefficients take only finitely many values is either rational or transcendental [8]. However, if the rational power series $A(X) = \sum_{n\geq 0} a_n X^n$ has bounded integer coefficients, then the sequence $(a_n)_{n\geq 0}$ must be ultimately periodic. Since the iPD sequence $\boldsymbol{u}$ is not ultimately periodic, we deduce that $U(X) = \sum_{n\geq 0} u_n X^n \in \mathbb{C}[[X]]$ is transcendental over $\mathbb{C}(X)$. □

5 Characteristic sequence of 1's in the iPD sequence $\boldsymbol{u}$

In this section, we study the characteristic sequence of 1's in the iPD sequence $\boldsymbol{u}$. The main result is that this sequence is not k-regular for any $k \geq 2$. Surprisingly, it is related to the characteristic sequence of Fibonacci numbers.

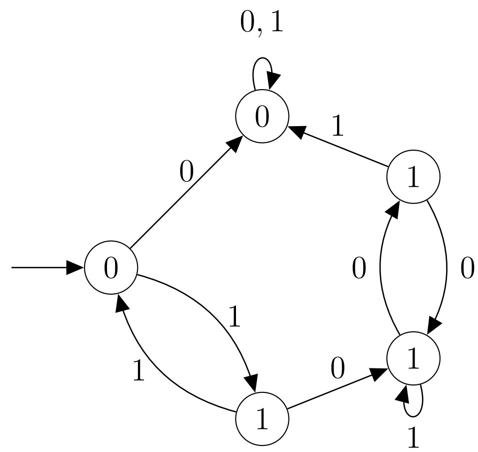

Figure 2: The 2-DFAO generating the inverse period-doubling sequence $\boldsymbol{u}$.

Definition 15. Let us define an increasing sequence $\boldsymbol{a} = (a_n)_{n\ge 0}$ satisfying $\{a_n \mid n \in N\} = \{m \in N \mid u_m = 1\}$ (sequence `A317543` in [15]). We have

$$\boldsymbol{a} = 1, 5, 7, 13, 17, 23, 29, 31, 37, 49, 55, 61, 65, 71, 77, 95, 101, 113, 119, 125, 127, 133, 145, \ldots.$$

From Proposition 10, we already know that $\boldsymbol{a}$ only contains odd integers. In the 2-DFAO in Figure 2, if the states outputting 1 are considered to be final, then the binary expansions of the terms of $\boldsymbol{a}$ is the language

$$L_a = \{\mathrm{rep}_2(a_n) \mid n \ge 0\} = \{11\}^*1 \cup 1\{1, 00\}^*0\{11\}^*1.$$

For instance, $\mathrm{rep}_2(a_0) = 1$, $\mathrm{rep}_2(a_1) = 101$, $\mathrm{rep}_2(a_2) = 111$, $\mathrm{rep}_2(a_3) = 1101$.

In the following, we obtain the complexity function of the language L_a. To that aim, we define the sequence $(F(n))_{n\ge 0}$ of the Fibonacci numbers with initial conditions equal to 0 and 1, i.e., $F(0) = 0$, $F(1) = 1$ and, for all $n \ge 2$, let $F(n) = F(n-1) + F(n-2)$.

Proposition 16. *The complexity function $\rho_{L_a} : \mathbb{N} \to \mathbb{N}$ of the language L_a satisfies $\rho_{L_a}(0) = 0$, $\rho_{L_a}(2n) = F(2n-1) - 1$ for all $n \ge 1$, and $\rho_{L_a}(2n+1) = F(2n) + 1$ for all $n \ge 0$.*

Proof. Start with the 2-DFAO in Figure 2. If the states outputting 1 are final, then it accepts the language L_a if words are read starting with the least significant digit. Let us reverse this 2-DFAO to obtain the non-deterministic automaton in Figure 3. It accepts L_a if words are read starting with the most significant digit. Determinizing it leads to an automaton $\mathcal{A}_{L_a}$ whose state set is $\{1, \ldots, 8\}$, with 1 as initial state, with 2, 6 as final states, and whose transition table is given in Table 1. Observe that this automaton accepts L_a, and has 8 as a sink state.

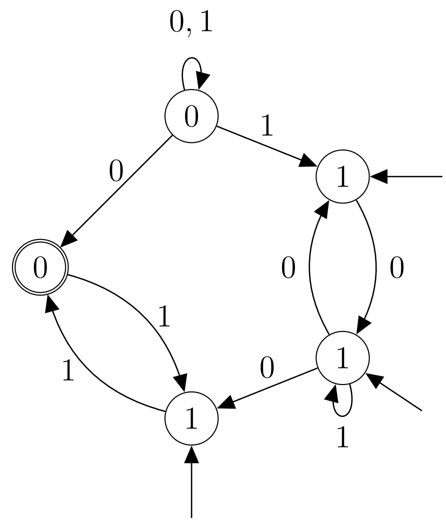

Figure 3: A non-deterministic automaton accepting the language L_a.

	0	1
1	8	2
2	3	4
3	5	6
4	3	2
5	3	5
6	8	7
7	8	6
8	8	8

Table 1: The transition table of a deterministic automaton accepting L_a.

Now let M denote the adjacency matrix of $\mathcal{A}_{L_a}$ without considering the sink state 8, and let

$$u = (1,0,0,0,0,0,0) \text{ and } v = (0,1,0,0,0,1,0)^\top.$$

Note that v codes the final states of $\mathcal{A}_{L_a}$. For all $n \geq 0$, $uM^n v$ is the number of length-n words in L_a. An easy induction shows that

$$M^{2n}v = (F(2n-1)-1, F(2n)+1, F(2n-1)-1, F(2n), F(2n), 1, 0)^\top \quad \forall n \geq 1,$$

and

$$M^{2n+1}v = (F(2n)+1, F(2n+1)-1, F(2n)+1, F(2n+1), F(2n+1)-1, 0, 1)^\top \quad \forall n \geq 0.$$

The conclusion follows. $\square$

The sequence $(a_n \bmod 3)_{n\geq 0}$ shows a particularly unexpected behavior as explained in the next two results.

Lemma 17. *Let $n \geq 0$. Then $a_n \bmod 3 \equiv r$ with $r \in \{1,2\}$. More precisely, let $w_n := \mathrm{rep}_2(a_n)$, and define $L_{a,1} = \{11\}^*1$ and $L_{a,2} = 1\{1,00\}^*0\{11\}^*1$. If $w_n \in L_{a,1}$, or if $w_n \in L_{a,2}$ and $|w_n|$ is even, then $a_n \bmod 3 \equiv 1$; if $w_n \in L_{a,2}$ and $|w_n|$ is odd, then $a_n \bmod 3 \equiv 2$.*

Proof. First, we have

$$(2^n \bmod 3)_{n\geq 0} = (1, -1, 1, -1, 1, -1, \ldots). \tag{8}$$

Now let $n \geq 0$ and set $w_n := \mathrm{rep}_2(a_n)$. If $w_n \in L_{a,1}$, then from (8) we deduce that $a_n \bmod 3 \equiv 1$. Assume that $w_n \in L_{a,2}$ and write $w_n = p_n s_n$ with $p_n \in 1\{1,00\}^*$ and $s_n \in 0\{11\}^*1$. Since $|s_n|$ is even, then (8) shows that $\mathrm{val}_2(s_n) \bmod 3 \equiv 1$.

As a first case, suppose that $|w_n|$ is odd. Then $|p_n|$ is also odd, and so p_n contains an odd number of 1's separated by even-length blocks of 0's. Because the 0's blocks have even length, the contributions of successive 1's in p_n alternate in value between $+1 \bmod 3$ and $-1 \bmod 3$. Since $|s_n|$ is even, after reading s_n then reading p_n gives an additional $+1 \bmod 3$. Consequently, both p_n and s_n together give $2 \bmod 3$, i.e., $a_n \bmod 3 \equiv \mathrm{val}_2(p_n s_n) \bmod 3 \equiv 2$.

As a second case, assume that $|w_n|$ is even. Then $|p_n|$ is even, and so p_n contains an even number of 1's separated by even-length blocks of 0's. Again the 1's in p_n contribute alternating $+1 \bmod 3$ and $-1 \bmod 3$, and since there is an even number of them, the 1's in p_n contribute $0 \bmod 3$ in total. Thus, in this case, $a_n \bmod 3 \equiv \mathrm{val}_2(p_n s_n) \bmod 3 \equiv 1$. ☐

Proposition 18. *The sequence $(a_n \bmod 3)_{n\geq 0}$ is given by the infinite word*

$$1^{F(1)}2^{F(2)}1^{F(3)}2^{F(4)}1^{F(5)}2^{F(6)}\cdots.$$

In particular, the sequence of run lengths of $(a_n \bmod 3)_{n\geq 0}$ is the sequence of Fibonacci numbers $(F(n))_{n\geq 1}$.

Proof. Recall that $L_a^n = L_a \cap \{0,1\}^n$ denotes the set of length-n words in L_a. We can order the words of L_a^n by lexicographic order, i.e.,

$$L_a^n = \{w_{n,1} <_{\mathrm{lex}} w_{n,2} <_{\mathrm{lex}} \cdots <_{\mathrm{lex}} w_{n,\#L_a^n}\}.$$

By Proposition 16, $\#L_a^0 = 0 = \#L_a^2$, $\#L_a^{2n} = F(2n-1) - 1$ for all $n \geq 1$, and $\#L_a^{2n+1} = F(2n) + 1$ for all $n \geq 0$.

Let us first consider L_a^{2n} for $n \geq 2$. From Lemma 17, we know that $\mathrm{val}_2(w_{2n,i}) \bmod 3 \equiv 1$ for all $i \in \{1, 2, \ldots, F(2n-1) - 1\}$. In other terms, we get

$$(\mathrm{val}_2(w_{2n,i}) \bmod 3)_{1\leq i\leq F(2n-1)-1} = 1^{F(2n-1)-1}.$$

Let us now study L_a^{2n+1} for $n \geq 0$. In the case where $n = 0$, then $L_a^1 = \{w_{1,1}\}$ with $w_{1,1} = 1$, which of course gives $\mathrm{val}_2(w_{1,1}) \bmod 3 = 1^{F(1)}$. Assume that $n \geq 1$. Since the words of

L_a^{2n+1} are ordered lexicographically, we know that $w_{2n+1,i} \in L_{a,2}$ for all $i \in \{1, 2, \ldots, F(2n)\}$, and $w_{2n+1,F(2n)+1} = 1^{2n+1} \in L_{a,1}$. From Lemma 17, we obtain that $\mathrm{val}_2(w_{2n+1,i}) \bmod 3 \equiv 2$ for all $i \in \{1, 2, \ldots, F(2n)\}$, and $\mathrm{val}_2(w_{2n+1,F(2n)+1}) \bmod 3 \equiv 1$. In fact, we obtain

$$(\mathrm{val}_2(w_{2n+1,i}) \bmod 3)_{1 \le i \le F(2n)+1} = 2^{F(2n)} 1.$$

Observe that, for any $n \ge 1$, concatenating the sequences $(\mathrm{val}_2(w_{2n+1,i}) \bmod 3)_{1 \le i \le F(2n)+1}$ and $(\mathrm{val}_2(w_{2n+2,i}) \bmod 3)_{1 \le i \le F(2n+1)-1}$ gives $(2^{F(2n)}1) \cdot (1^{F(2n+1)-1}) = 2^{F(2n)}1^{F(2n+1)}$. Now putting everything together, we find

$$\begin{aligned}(a_n \bmod 3)_{n \ge 0} &= 1^{F(1)} \cdot 2^{F(2)}1 \cdot 1^{F(3)-1} \cdot 2^{F(4)}1 \cdot 1^{F(5)-1} \cdot 2^{F(6)}1 \cdots \\ &= 1^{F(1)}2^{F(2)}1^{F(3)}2^{F(4)}1^{F(5)}2^{F(6)} \cdots,\end{aligned}$$

as expected. ☐

To show that $\boldsymbol{a}$ is not k-regular for any $k \ge 2$, the idea is to study the sequence of consecutive differences in $(a_n \bmod 3)_{n \ge 0}$. Let us define the sequence $\boldsymbol{\delta} = (\delta_n)_{n \ge 0}$ by

$$\delta_n = \begin{cases} 1, & \text{if } (a_{n+1} - a_n) \bmod 3 \neq 0; \\ 0, & \text{otherwise.} \end{cases}$$

From Proposition 18, we know that $\delta_n = 1$ if and only if there exists $n = F(m) - 2$ for some $m \ge 0$. If we let $\boldsymbol{x}$ denote the characteristic sequence of Fibonacci numbers, i.e., x_n equals 1 if n is a Fibonacci number, 0 otherwise, then $\boldsymbol{\delta} = (x_n)_{n \ge 2}$ since for all $n \ge 0$

$$\delta_n = 1 \Leftrightarrow n = F(m) - 2 \text{ for some } m \ge 0 \Leftrightarrow n + 2 = F(m) \text{ for some } m \ge 0 \Leftrightarrow x_{n+2} = 1.$$

The goal is now to show that $\boldsymbol{x}$ is not k-automatic for any $k \ge 2$; then the non-k-automaticity of $\boldsymbol{\delta}$ can easily be deduced. What follows is widely inspired by [12, 13]. In our context, we consider the ANS $(L_F, \{0, 1\}, <)$ where $L_F = \{\varepsilon\} \cup 1\{0, 01\}^*$ is the language of Fibonacci representations of nonnegative integers with $0 < 1$. Observe that the DFA $\mathcal{A}$ in Figure 4 accepts the regular language L_F.

Lemma 19. *The characteristic sequence of Fibonacci numbers $\boldsymbol{x}$ is Fibonacci-automatic.*

Proof. The Fibonacci-DFAO $\mathcal{B}$ in Figure 5 generates the sequence $\boldsymbol{x}$ in the Zeckendorff numeration system. In particular, this shows that $\boldsymbol{x}$ is Fibonacci-automatic. ☐

When a word is S-automatic for some ANS S, then it is in fact morphic [13].

Theorem 20. *An infinite word $\boldsymbol{w}$ is morphic if and only if $\boldsymbol{w}$ is S-automatic for some ANS S.*

From Lemma 19 and Theorem 20, we easily deduce that $\boldsymbol{x}$ is morphic. More precisely, we want to build the morphisms that generate $\boldsymbol{x}$. We follow the constructive proof of Theorem 20 (we refer the reader to [13, Chapter 2] for more details).

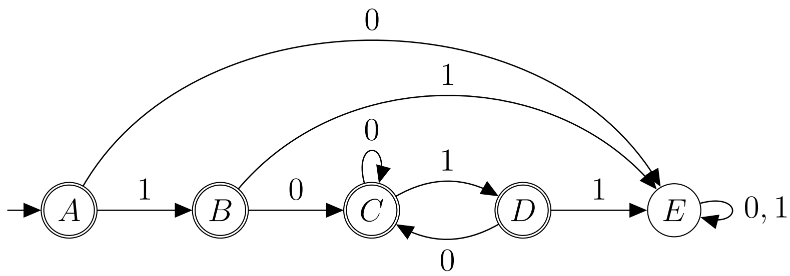

Figure 4: The DFA $\mathcal{A}$ accepting the language $\{\varepsilon\} \cup 1\{0, 01\}^*$.

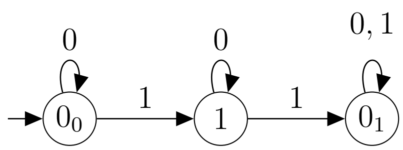

Figure 5: The Fibonacci-DFAO $\mathcal{B}$ generating $\boldsymbol{x}$.

Lemma 21. *Let $f : \{z, a_0, a_1, \ldots, a_7\}^* \to \{z, a_0, a_1, \ldots, a_7\}^*$ be the morphism defined by $f(z) = za_0$ and*

i	0	1	2	3	4	5	6	7
$f(a_i)$	a_1a_2	a_1a_4	a_3a_7	a_3a_6	a_4a_7	a_5a_6	a_5a_7	a_7a_7

We also define the morphism $g : \{z, a_0, a_1, \ldots, a_7\}^* \to \{0, 1\}^*$ by $g(z) = g(a_1) = g(a_4) = g(a_7) = \varepsilon$, $g(a_0) = g(a_5) = g(a_6) = 0$ and $g(a_2) = g(a_3) = 1$. Then $\boldsymbol{x} = g(f^\omega(z))$. In particular, the word $\boldsymbol{x}$ is morphic.

Proof. First recall that the DFA $\mathcal{A}$ in Figure 4 accepts the language $L_F = \{\varepsilon\} \cup 1\{0, 01\}^*$, and the Fibonacci-DFAO $\mathcal{B}$ in Figure 5 generates the sequence $\boldsymbol{x}$. Then, the product automaton $\mathcal{P} = \mathcal{A} \times \mathcal{B}$ is drawn in Figure 6. If we set

$$\begin{aligned} &a_0 := (A, 0_0), a_1 := (E, 0_0), a_2 := (B, 1), a_3 := (C, 1), \\ &a_4 := (E, 1), a_5 := (C, 0_1), a_6 := (D, 0_1), a_7 := (E, 0_1), \end{aligned}$$

then we can associate a morphism $\psi_{\mathcal{P}} : \{z, a_0, a_1, \ldots, a_7\}^* \to \{z, a_0, a_1, \ldots, a_7\}^*$ with $\mathcal{P}$ as follows. It is defined by $\psi_{\mathcal{P}}(z) = za_0$ and

i	0	1	2	3	4	5	6	7
$\psi_{\mathcal{P}}(a_i) = \delta_{\mathcal{P}}(a_i, 0)\delta_{\mathcal{P}}(a_i, 1)$	a_1a_2	a_1a_4	a_3a_7	a_3a_6	a_4a_7	a_5a_6	a_5a_7	a_7a_7

where $\delta_{\mathcal{P}}$ is the transition function of $\mathcal{P}$. Notice that $\psi_{\mathcal{P}} = f$. We also define the morphism

$$g : \{z, a_0, a_1, \ldots, a_7\}^* \to \{0, 1\}^* : z, a_1, a_4, a_7 \mapsto \varepsilon; a_0, a_5, a_6 \mapsto 0; a_2, a_3 \mapsto 1.$$

It is well known that $\boldsymbol{x} = g(f^\omega(z))$, which shows that $\boldsymbol{x}$ is morphic. $\square$

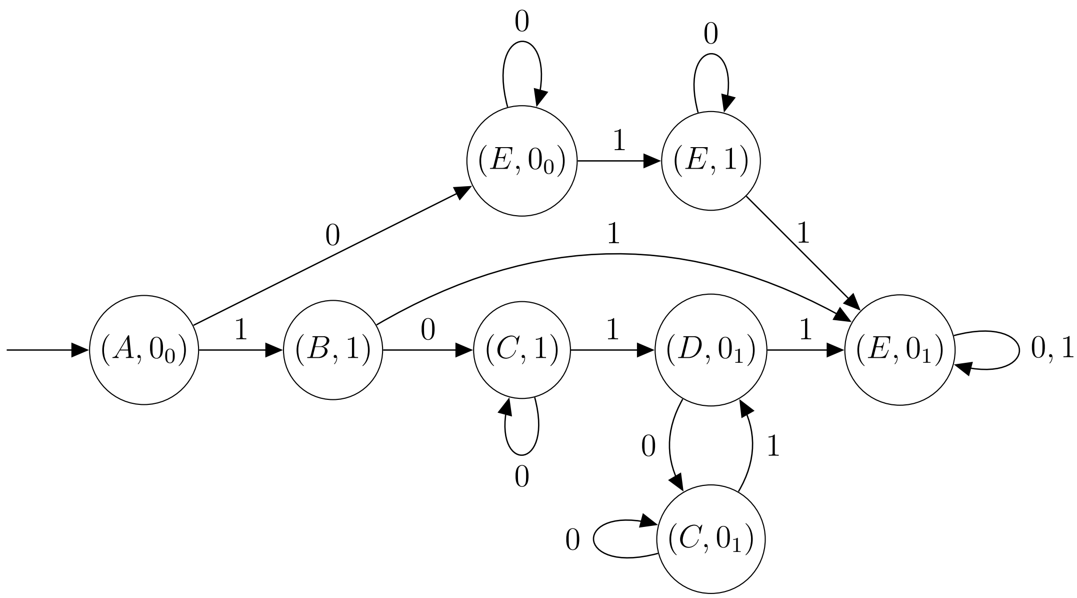

Figure 6: The DFA $\mathcal{P}$ which is the product of $\mathcal{A}$ and $\mathcal{B}$.

Observe that the morphism g in Lemma 21 is erasing, i.e., the image of some letter is the empty word. In the following lemma (see [12, Chapter 3]), we get rid of the erasure and we later obtain two new non-erasing morphisms that generate $\boldsymbol{x}$.

Lemma 22. *Let $\boldsymbol{w} = g(f^\omega(a))$ be a morphic word where $g : B^* \to A^*$ is a (possibly erasing) morphism and $f : B^* \to B^*$ is a non-erasing morphism. Let C be a subalphabet of $\{b \in B \mid g(b) = \varepsilon\}$ such that f_C is a submorphism of f. Let $\lambda_C : B^* \to B^*$ be the morphism defined by $\lambda_C(b) = \varepsilon$ if $b \in C$, and $\lambda_C(b) = b$ otherwise. The morphisms $f_\varepsilon := (\lambda_C \circ f)|_{(B\setminus C)^*}$ and $g_\varepsilon := g|_{(B\setminus C)^*}$ are such that $\boldsymbol{w} = g_\varepsilon(f_\varepsilon^\omega(a))$.*

Proposition 23. *Let $\phi : \{a,b,c,d,e\}^* \to \{a,b,c,d,e\}^*$ be the morphism defined by*

$$\phi : \{a,b,c,d,e\}^* \to \{a,b,c,d,e\}^* : \begin{cases} a & \mapsto & ab, \\ b & \mapsto & c, \\ c & \mapsto & ce, \\ d & \mapsto & de, \\ e & \mapsto & d \end{cases}$$

and let $\mu : \{a,b,c,d,e\}^* \to \{0,1\}^* : a,d,e \mapsto 0; b,c \mapsto 1$ be a coding. Then $\boldsymbol{x} = \mu(\phi^\omega(a))$.

Proof. We make use of Lemmas 21 and 22. First, we have

$$\{b \in \{z, a_0, a_1, \ldots, a_7\} \mid g(b) = \varepsilon\} = \{z, a_1, a_4, a_7\},$$

so we choose $C = \{a_1, a_4, a_7\}$ for f_C is a submorphism of f. Then the morphism

$$f_\varepsilon : \{z, a_0, a_2, a_3, a_5, a_6\}^* \to \{z, a_0, a_2, a_3, a_5, a_6\}^*$$

is defined by $f_\varepsilon(z) = za_0$, $f_\varepsilon(a_0) = a_2$, $f_\varepsilon(a_2) = a_3$, $f_\varepsilon(a_3) = a_3a_6$, $f_\varepsilon(a_5) = a_5a_6$ and $f_\varepsilon(a_6) = a_5$, while the morphism $g_\varepsilon : \{z, a_0, a_2, a_3, a_5, a_6\}^* \to \{0,1\}^*$ is given by $g_\varepsilon(z) = \varepsilon$, $g_\varepsilon(a_0) = g_\varepsilon(a_5) = g_\varepsilon(a_6) = 0$ and $g_\varepsilon(a_2) = g_\varepsilon(a_3) = 1$. We also have $\boldsymbol{x} = g_\varepsilon(f_\varepsilon^\omega(z))$. Note that $f_\varepsilon|_{\{a_2,a_3,a_5,a_6\}^*}$ is a submorphism of f_ε.

Let us define the morphism $f'_\varepsilon : \{a_0, a_2, a_3, a_5, a_6\}^* \to \{a_0, a_2, a_3, a_5, a_6\}^*$ by $f'_\varepsilon(a_0) = a_0a_2$, and $f'_\varepsilon = f_\varepsilon|_{\{a_2,a_3,a_5,a_6\}^*}$. From that definition, f'_ε is prolongable on a_0. Also consider the morphism $g'_\varepsilon : \{a_0, a_2, a_3, a_5, a_6\}^* \to \{0,1\}^*$ given by $g'_\varepsilon = g_\varepsilon|_{\{a_0,a_2,a_3,a_5,a_6\}^*}$. We have

$$\begin{aligned}
f_\varepsilon^\omega(z) &= za_0 f_\varepsilon(a_0) f_\varepsilon^2(a_0) f_\varepsilon^3(a_0) f_\varepsilon^4(a_0)\cdots \\
&= za_0 f_\varepsilon(a_0) f_\varepsilon(f_\varepsilon(a_0)) f_\varepsilon^2(f_\varepsilon(a_0)) f_\varepsilon^3(f_\varepsilon(a_0))\cdots \\
&= za_0a_2 f_\varepsilon(a_2) f_\varepsilon^2(a_2) f_\varepsilon^3(a_2)\cdots \\
&= za_0a_2 f'_\varepsilon(a_2)(f'_\varepsilon(a_2))^2(f'_\varepsilon(a_2))^3\cdots,
\end{aligned}$$

thus we get

$$\begin{aligned}
\boldsymbol{x} &= g_\varepsilon(f_\varepsilon^\omega(z)) \\
&= g_\varepsilon(z)g_\varepsilon(a_0)g_\varepsilon(a_2)g_\varepsilon(f'_\varepsilon(a_2))g_\varepsilon((f'_\varepsilon(a_2))^2)g_\varepsilon((f'_\varepsilon(a_2))^3)\cdots \\
&= \varepsilon g'_\varepsilon(a_0)g'_\varepsilon(a_2)g'_\varepsilon(f'_\varepsilon(a_2))g'_\varepsilon((f'_\varepsilon(a_2))^2)g'_\varepsilon((f'_\varepsilon(a_2))^3)\cdots \\
&= g'_\varepsilon(a_0a_2f'_\varepsilon(a_2)(f'_\varepsilon(a_2))^2(f'_\varepsilon(a_2))^3\cdots) \\
&= g'_\varepsilon((f'_\varepsilon)^\omega(a_0)).
\end{aligned}$$

Up to a renaming of the letters, we have proven the claim. ☐

Corollary 24. *Let $\varphi = \frac{1}{2}(\sqrt{5}+1)$ be the golden ratio. The word $\boldsymbol{x}$ is φ-substitutive.*

Proof. Let

$$M_\phi = \begin{pmatrix} 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

be the matrix associated with the morphism ϕ. The Perron-Frobenius eigenvalue of M_ϕ is $\varphi = \frac{1}{2}(\sqrt{5}+1)$. Since all the letters of $\{a, b, c, d, e\}$ occur in $\phi^\omega(a)$, then $\boldsymbol{x}$ is φ-substitutive by Proposition 23. ☐

Proposition 25. *The sequence $\boldsymbol{x}$ is not k-automatic for any $k \in \mathbb{N}_{\geq 2}$.*

Proof. Proceed by contradiction and suppose that there exists an integer $k \geq 2$ such that $\boldsymbol{x}$ is k-automatic. Then, by Theorem 2, $\boldsymbol{x}$ is also k-substitutive. Indeed, it is not difficult to see that the Perron-Frobenius eigenvalue of the matrix associated with a k-uniform morphism is the integer k. Clearly, k and φ are two multiplicatively independent real numbers. Thus, by Theorem 1, $\boldsymbol{x}$ is ultimately periodic. This is impossible. ☐

Corollary 26. *The sequence $(a_n)_{n\geq 0}$ is not k-regular for any $k \in \mathbb{N}_{\geq 2}$.*

Proof. Suppose that the sequence $(a_n)_{n\geq 0}$ is k-regular for some $k \geq 2$. Then by Proposition 3, the sequence $(a_n \bmod 3)_{n\geq 0}$ is k-automatic, and so is $\boldsymbol{x}$. This contradicts Proposition 25. $\square$

We end this section with the following open problem.

Problem 27. Let us define an increasing sequence $\boldsymbol{b} = (b_n)_{n\geq 0}$ satisfying $\{b_n \mid n \in N\} = \{m \in N \mid u_m = 0\}$ (sequence `A317544` in [15]). This is the characteristic sequence of 0's in the iPD sequence $\boldsymbol{u}$. We have

$$\boldsymbol{b} = 0, 2, 3, 4, 6, 8, 9, 10, 11, 12, 14, 15, 16, 18, 19, 20, 21, 22, 24, 25, 26, 27, 28, 30, 32, 33, 34, 35, \ldots.$$

Is the sequence $\boldsymbol{b}$ k-regular for some $k \geq 2$?

Remark 28. After submission of the present paper for publication, Ł. Merta contacted us with a different proof of Corollary 26 that uses his work on the Baum-Sweet sequence [10]. In this note, we only sketch this other proof. We let $\boldsymbol{bs} = (bs_n)_{n\geq 0}$ denote the Baum-Sweet sequence (indexed by `A086747` in [15]). First, using Proposition 10 and recurrence relations satisfied by $\boldsymbol{bs}$, one can prove that $u_{2n+1} = bs_{n+1}$ for all $n \geq 0$. Then, if $\boldsymbol{\ell} = (\ell_n)_{n\geq 0}$ is the characteristic sequence of 1's in the Baum-Sweet sequence $\boldsymbol{bs}$, we deduce that $a_n = 2\ell_{n+1} - 1$ for all $n \geq 0$. From the stability properties of k-regular sequences (see [4]), if $\boldsymbol{a}$ was k-regular for some $k \geq 2$, then $\boldsymbol{\ell}$ would also be, which contradicts [10, Theorem 2.6].

6 Acknowledgment

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